Mathematical Model of Computer Viruses



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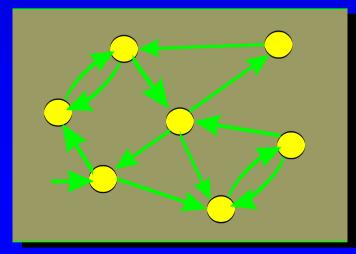


- Models of computation
- Operating system
- Virus definition
- What can we do with this mathematical model?



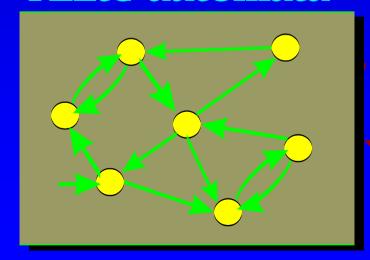


Finite automata

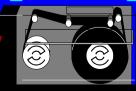




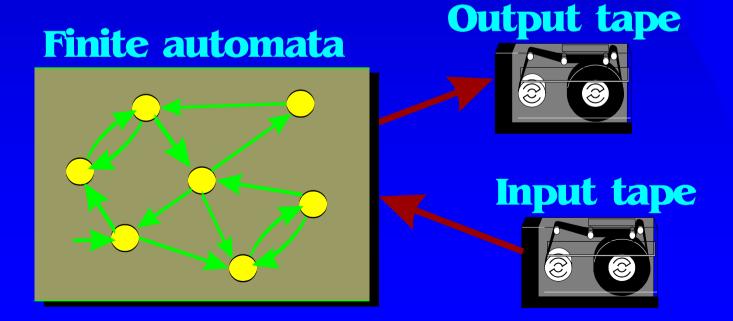
Finite automata



Input tape







 $T = \langle Q, S, I, \delta, b, q_0, q_f \rangle$



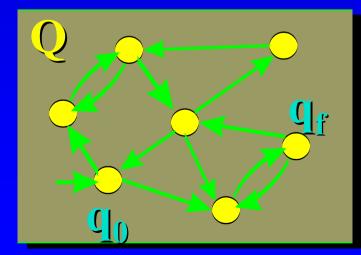
S: tape symbols

I: input symbols, $I \subset S$

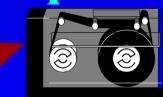
b: blank symbol, $b \in S \setminus I$

δ: move function, $\delta: \mathbf{Q} \times \mathbf{S} \to \mathbf{Q} \times \mathbf{S} \times \{\mathbf{l}, \mathbf{r}, \mathbf{s}\}$

Finite automata

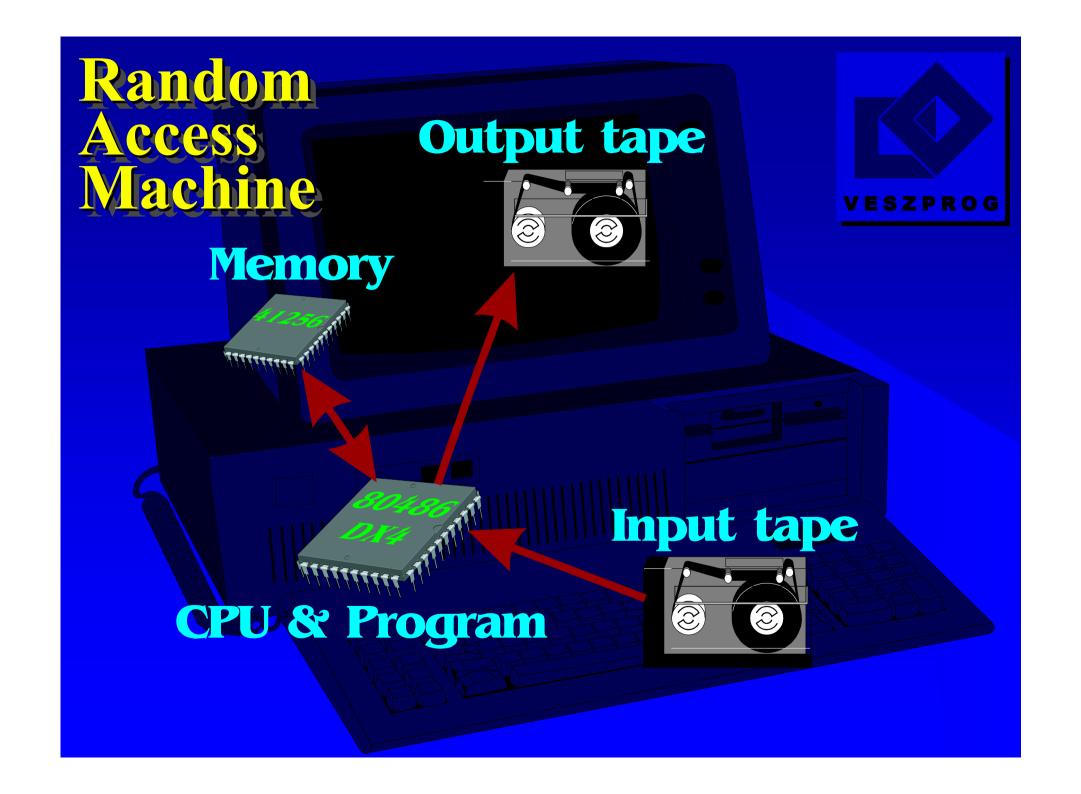


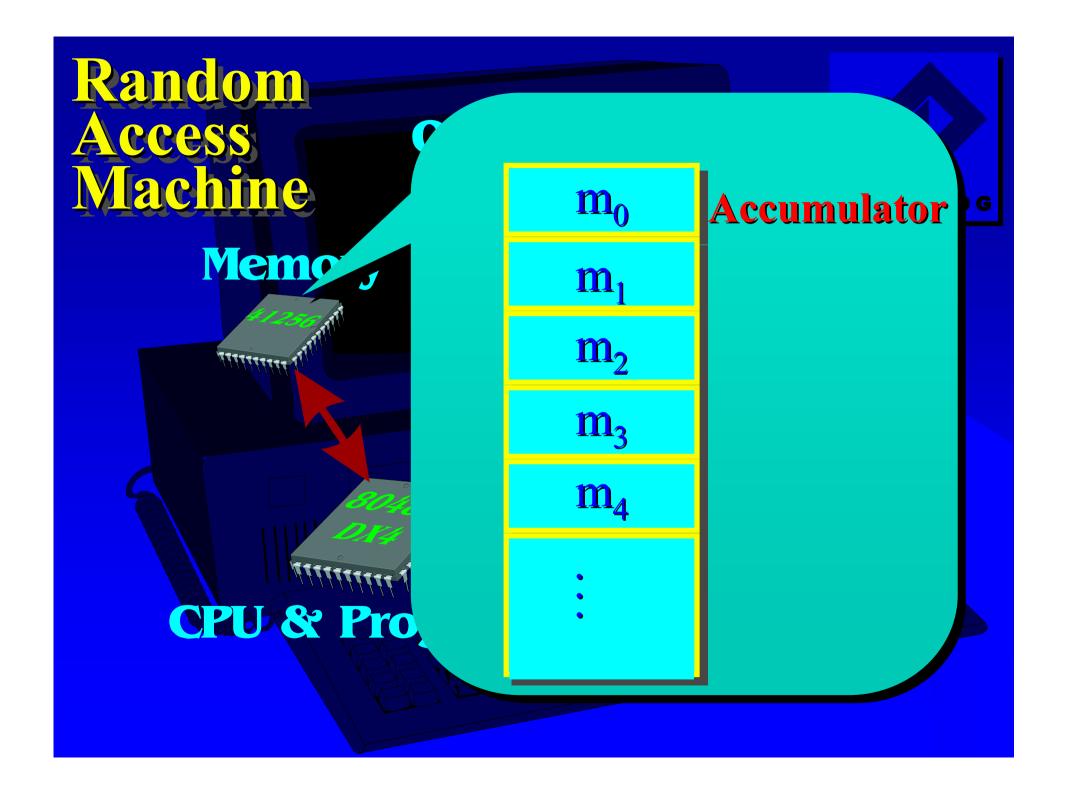
Output tape

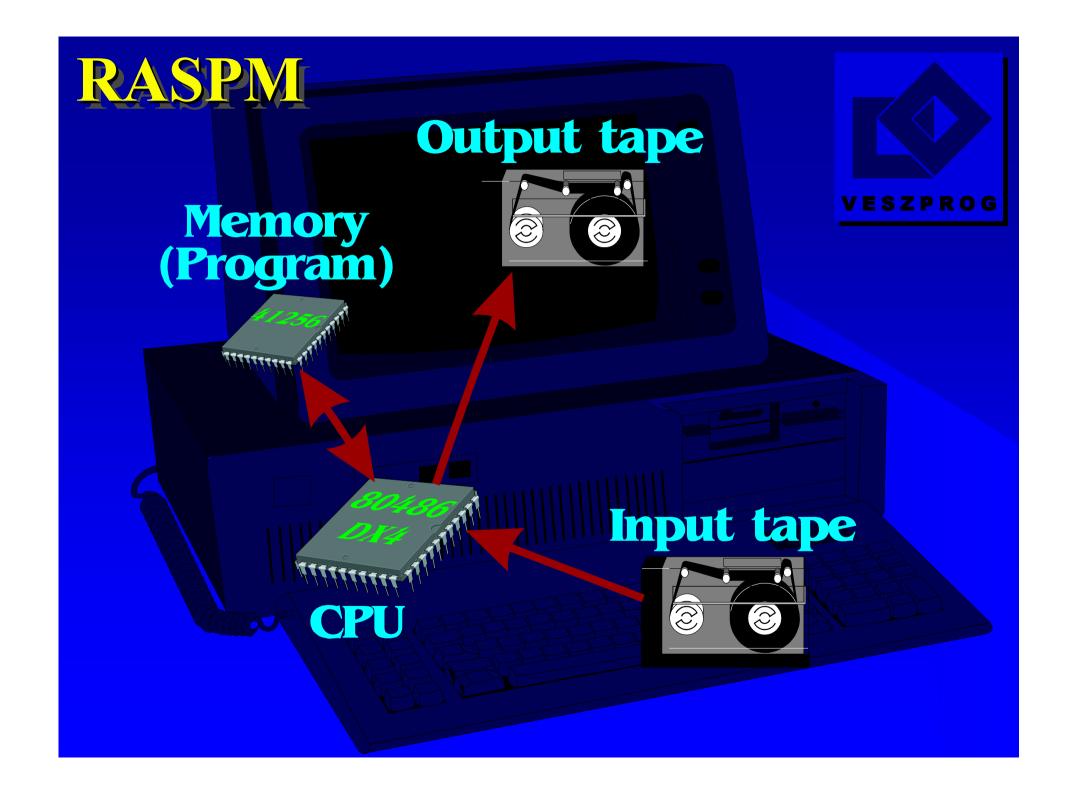


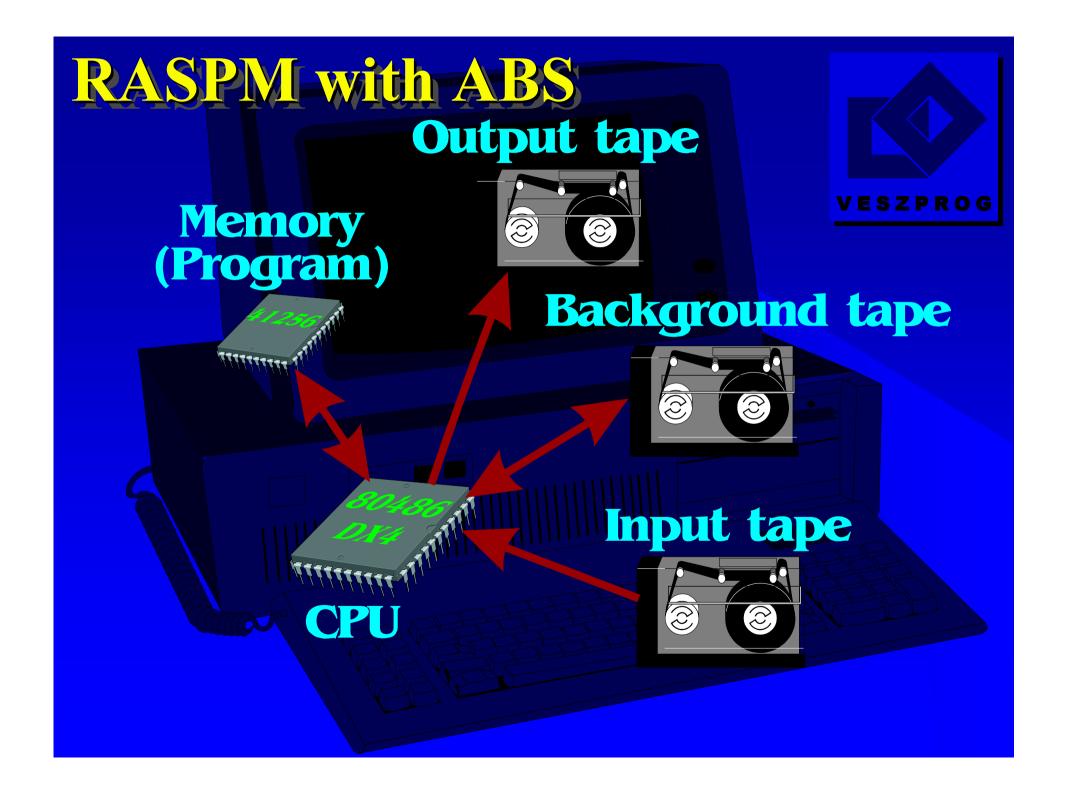
Input tape

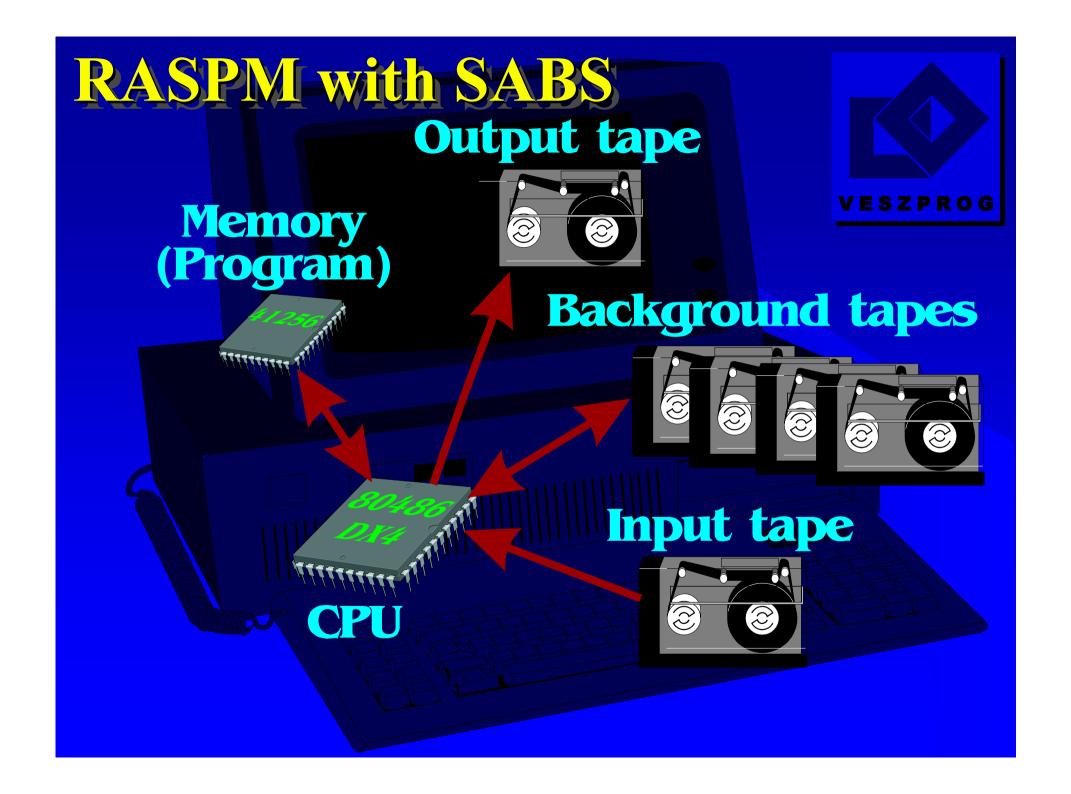






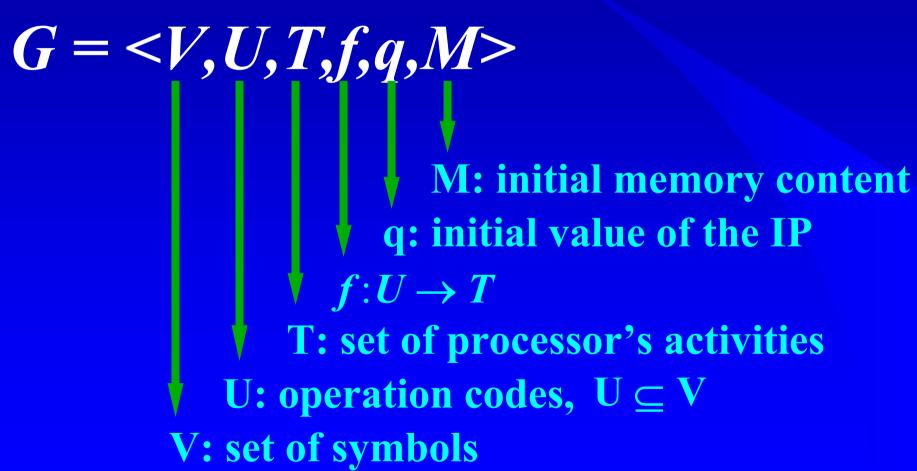






RASPM with ABS definition





Instruction set



- move (LOAD, STORE)
- logical (AND, OR, XOR)
- arithmetic (ADD, SUB, MULT, DIV)
- branch (JUMP, JGTZ, JZERO)
- input/output tape handling (READ, WRITE)
- background tape handling (GET, PUT, SEEK, SETDRIVE)





system of programs



- system of programs
- able to handle separate program or data files



- system of programs
- able to handle separate program or data files
- able to make a specified program to run.





• The OS is in the initial memory (M)



- The OS is in the initial memory (M)
 - → OS specific machine



- The OS is in the initial memory (M)
 → OS specific machine
- The OS is in the background tape



- The OS is in the initial memory (M)
 - → OS specific machine
- The OS is in the background tape
 - → OS independent machine



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- The OS is in the initial memory (M)
 - → OS specific machine
- The OS is in the background tape
 - → OS independent machine
- The OS is in the input tape
 - → unusable



- The OS is in the initial memory (M)
 → OS specific machine
- The OS is in the background tape
 → OS independent machine
- The OS is in the input tape
 → unusable

Sample OS





$$G_1 = \langle V_1, U_1, T_1, f_1, q_1, M_1 \rangle$$

 $G_2 = \langle V_2, U_2, T_2, f_2, q_2, M_2 \rangle$



$$G_1 = \langle V_1, U_1, T_1, f_1, q_1, M_1 \rangle$$

 $G_2 = \langle V_2, U_2, T_2, f_2, q_2, M_2 \rangle$
 $\{q_1, M_1\} \neq \{q_2, M_2\}$



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 $\{q_1, M_1\} \neq \{q_2, M_2\}$

- different operating systems
- different loader program



$$G_1 = \langle V_1, U_1, T_1, f_1, q_1, M_1 \rangle$$

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- different activities
- different operation codes



$$G_1 = \langle V_1, U_1, T_1, f_1, q_1, M_1 \rangle$$

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$$V_1 \neq V_2$$



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$$V_1 \neq V_2$$

- different symbols
- different tape formats





a (part of) program



- a (part of) program
- it is attached to a program area



- a (part of) program
- it is attached to a program area
- it is able to link itself to other program areas



- a (part of) program
- it is attached to a program area
- it is able to link itself to other program areas
- it is executed when the host program area is to be executed





machine specific



- machine specific
- machine independent



- machine specific
- machine independent
- operating system specific



- machine specific
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- machine specific
- machine independent
- operating system specific
- operating system independent
- direct



- machine specific
- machine independent
- operating system specific
- operating system independent
- direct
- indirect



- machine specific
- machine independent
- operating system specific
- operating system independent
- direct
- indirect

Sample virus







• Examine the working mechanism of viruses



- Examine the working mechanism of viruses
- Examine the virus detection problem



- Examine the working mechanism of viruses
- Examine the virus detection problem
- Examine multiplatform viruses



Theorem:

It is impossible to build a Turing Machine which could decide if an executable file in a RASPM with ABS contains a virus or not.



Proof:





Proof:

Host program Virus TM prg



Proof:

Host program Virus TM prg TM input

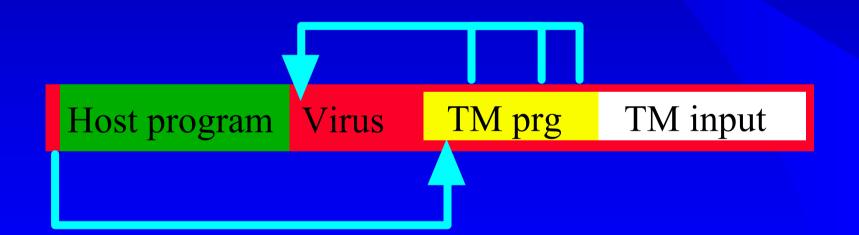


Proof:

Host program Virus TM prg TM input

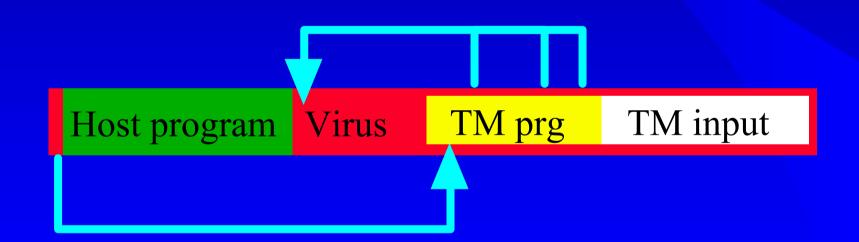


Proof:





Proof:



Virus detection problem — TM halting problem

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Conditions:

$$V_1 \stackrel{\text{\tiny W}}{\cup} U_2 \neq 0$$

$$U_1 \stackrel{\text{\tiny W}}{\cup} V_2 \neq 0$$

 G_1 has to know some operation codes of G_2 G_2 has to know some operation codes of G_1



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Conditions:

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- The virus code can be the same.



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$$U_1 \overset{\text{\tiny def}}{\cup} U_2 = 0$$

- The virus code must be different.

Future



Future



- Examine general virus detection problem in limited cases:
 - Spreading under the model
 - Limit the time/space

Future



- Examine general virus detection problem in limited cases:
 - Spreading under the model
 - Limit the time/space
- Examine polymorphic techniques
 - Without coding/decoding
 - Changing instructions









































• For what kind of viruses can be used?



- For what kind of viruses can be used?
- What is the probability of false alarms?



- For what kind of viruses can be used?
- What is the probability of false alarms?
- What is the expense criteria?

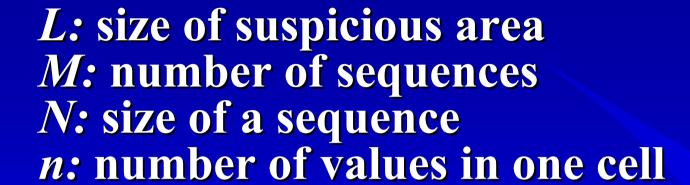
Sequence searching algorithm

VESZPROG

Sequence searching algorithm VESZEROG

• for non-polymorphic known viruses

Sequence searching algorithm





• for non-polymorphic known viruses

• false alarms:
$$p \approx \frac{L \cdot M}{n^N}$$

Sequence searching algorithm

L: size of suspicious areaM: number of sequencesN: size of a sequencen: number of values in one cell



for non-polymorphic known viruses

• false alarms:
$$p \approx \frac{L \cdot M}{n^N}$$

• expense criteria: P, polynomial $\leq L \cdot M \cdot N$ comparisions





for known viruses

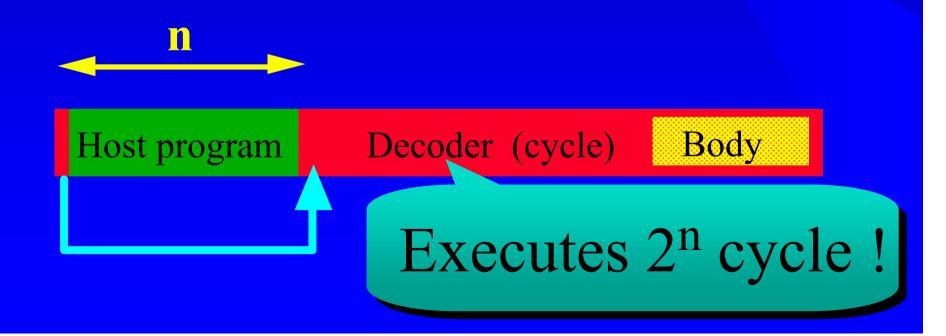


- for known viruses
- expense criteria:

Host program
Decoder (cycle)
Body



- for known viruses
- expense criteria: NP







Host program Decoder Body



Host program Decoder Body

$$\alpha = \frac{\text{size of variable parts of the virus}}{\text{full size of the virus}}$$



Host program Decoder Body

$$\alpha = \frac{\text{size of variable parts of the virus}}{\text{full size of the virus}}$$

 β = number of variants of the decoders





search for an uninfected program



search for an uninfected program

append virus



search for an uninfected program

append virus

choose a random instruction in the virus



search for an uninfected program

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choose a random instruction in the virus



search for an uninfected program

appen'd virus

choose a random instruction in the virus

repeat 100 times



search for an uninfected program

append virus

choose a random instruction in the virus

repeat 100 times



search for an uninfected program

append virus

choose a random instruction in the DISK

repeat 100 times

Name: RIPPER

Aliases: Jack Ripper

Status: Common

Origin: Norway

Length: 1024 bytes (2 sectors)

Infect: MBR, Boot sector

Other: Resident, Stealth,

Disk corruption



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The virus swaps two words in the DOS write buffer. It occurs on a random basis of approximately 1 write in 1024 cases.



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- Examine the working mechanism of viruses
- Examine the virus detection problem
- Examine multiplatform viruses
- Examine new polymorphic virus types